Baäerina Swan Lake

Code generation and intermediate representation in Ballerina Compiler

April 2024



Lecture Outline

- Evolution of Ballerina Compiler
- Current structure of Ballerina Compiler
- Intermediate representations
 - Types of representations
 - $\circ \quad \ \ \text{Lowering to IR}$

Evolution : WSO2's history in languages

WSO2 enables thousands of enterprises, including hundreds of the world's largest corporations, top universities, and governments, to drive their digital transformation journeys—executing more than 18 trillion transactions and managing more than 500 million identities annually.



Evolution of

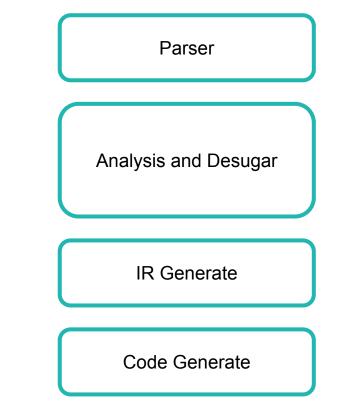


- Started out as Synapse replacement language back in late 2016. Inspired by sequence diagrams and graphical editing.
- Initial implementation as AST interpreted language (2017)
- Internal vm (BVM) with internal ByteCode (late 2017)
- Backend/frontend separation via BIR. JVM bytecode as the backend (late 2018).
- Swan Lake version GA release in , with major improvements and extensive set of standard libraries and connectors (early 2022).
- Continuous updates to Swan Lake version. Currently on update 8.5



Structure of







Features of Ballerina



Data oriented

Type-safe, declarative processing of JSON, XML, and tabular data with language-integrated queries.

```
type User record { int id; string name; };
...
User manu = { id: 92874, name: "manuranga" }
```



Concurrent

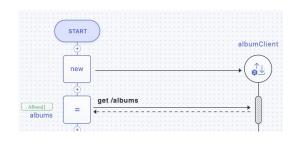
Easy and efficient concurrency with sequence diagrams and language-managed threads without the complexity of asynchronous functions.

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Graphical

Programs have both a textual syntax and an equivalent graphical form based on sequence diagrams. http:Client hello = check new ("http://hello.com");
MyGreeting greeting = check hello->get("/world");

Also see: start, wait and workers





Features of Ballerina



Flexibly typed

Uses structural types with support for openness for static typing within a program and for describing service interfaces.



Reliable, maintainable

Explicit error handling, static types, and concurrency safety, combined with a familiar, readable syntax make programs reliable and maintainable.



Cloud native

Network primitives in the language make it simpler to write services and run them in the cloud.

```
type Customer record {|
    int id;
    string name;
    int account;
|};
```

```
Customer customer = { ... };
User user = customer;
addUser(user);
```

> bal build Compiling source example/greeter:0.1.0

Generating executable

Generating artifacts...

@kubernetes:Service @kubernetes:Deployment @kubernetes:HPA @kubernetes:Docker

- complete 1/1 - complete 1/1 - complete 1/1 - complete 2/2



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Evolution : Parser

Parser Generator

```
functionSignature :
    LEFT_PARENTHESIS
    formalParameterList?
    RIGHT_PARENTHESIS
    returnParameter?;
```

Easier to get started Less boilerplate

We used ANTLR until late 2020

https://github.com/ballerina-platform/ballerina -lang/blob/v0.995.9/compiler/ballerina-lang/sr c/main/resources/grammar/BallerinaLexer.g4

Handwritten Recursive Descent

```
private STNode parseFuncSignature(boolean isParamNameOptional) {
   STNode openParenthesis = parseOpenParenthesis();
   STNode parameters = parseParamList(isParamNameOptional);
   STNode closeParenthesis = parseCloseParenthesis();
   endContext(); // end param-list
   STNode returnTypeDesc = parseFuncReturnTypeDescriptor(isParamNameOptional);
   return STNodeFactory.createFunctionSignatureNode(openParenthesis, parameters
   closeParenthesis, returnTypeDesc);
}
```

Faster for complex grammars Better at handling edge cases, more flexibility Better error recovery

It could be much simpler :-

https://eli.thegreenplace.net/2012/08/02/parsing-expressions-by-precedence-climbing

https://github.com/ballerina-platform/nballerina/blob/main/compiler/modules/front .syntax/parseExpr.bal



Evolution : Execution

AST Interpreter

public BValue visit(BinaryExpression binExpr) {
 Expression rExpr = binExpr.getRExpr();
 BValueType rValue = rExpr.execute(this);
 Expression lExpr = binExpr.getLExpr();
 BValueType lValue = lExpr.execute(this);
 return binExpr.getEvalFunc().apply(lValue, rValue);
}

Close to source language

ByteCode Interpreter

while (ip < code.length) { Instruction instruction = code[ip]; int[] operands = instruction.getOperands(); int opcode = instruction.getOpcode(); switch (opcode) { case InstructionCodes ICONST⁻ cpIndex = operands[0]: i = operands[1]; sf.longRegs[i] = ((IntegerCPEntry) constPool[cpIndex]).getValue(); break: case InstructionCodes ECONST¹ cpIndex = operands[0]; i = operands[1]; sf.doubleRegs[i] = ((FloatCPEntry) constPool[cpIndex]).getValue(); break;

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Reasonable speed (Order of Python) Used by Ballerina compiler until 2019

Compiled (to external format)

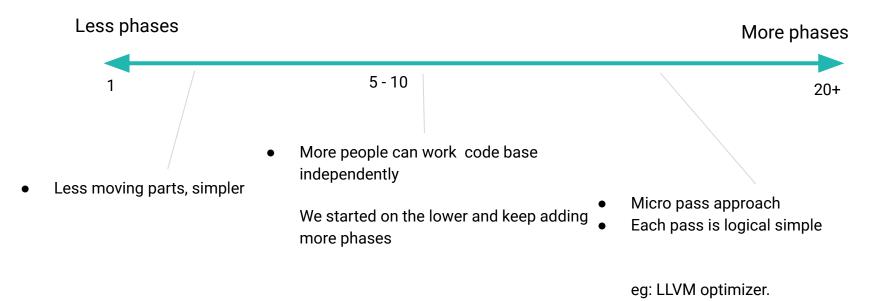
We are using JVM bytecode in Swan Lake Version

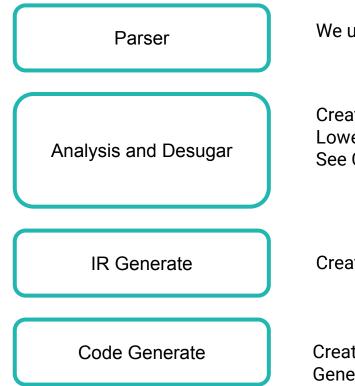
Experimented with LLVM and WebAssemblye



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Evolution : Compiler phases





We used to have ANTLR, now we have a custom parser

Create symbols, Type check, Run plugins Lower lambdas and other high level control flow See CompilerPhaseRunner Class of ballerina-lang repo

Create BIR. Conditionals get lowered to GOTOs

Create JVM Bytecode. Create concurrency yield points. Generate classes for values.





Intermediate representation

Q: How are compilers phases are connected? A: IR

Why IR (instead in-memory graph)

- Easy to debug due to serialization
- Can verify

Styles of IR

- Stack based vs Register based
- Flat vs structured
- SSA Register vs Mutable Register

IR Styles : Stack machine vs Register base

Stack machine

iload %1 iload 20 iadd

- May produce smaller IR.
- Used by JVM and WebAssembly

We didn't pick this option due to the added complexity in generation. a = b + 20;

Register base (mutable)

iload 20 %2 iadd %1 %2 %0

• More closer to source language.

We use this format in Ballerina (BVM, jBallerina IR) Register base (SSA) %a = add %b 20

- Better for analysis and optimization.
- Need phi nodes

We didn't pick because not a good input format for JVM (or for LLVM surprisingly, due to debug info)



IR Styles : Flat vs Structured

Flat IR

Structured IR

loopHead:

local.get \$i
i32.const 10
i32.gt_s loopEnd
goto loopHead

loopEnd:

- Most popular format
- Can result in non-reducible loops

(loop \$my_loop local.get \$i i32.const 10 i32.lt_s br_if \$my_loop)

- Used by WebAssembly
- No GOTOs

Ballerina IR Example

```
public function add(int a, int b) returns int {
    return a + b;
}
```

```
public add function(int, int) -> int {
    %0(RETURN) int;
    %1(ARG) int;
    %2(ARG) int;
    bb0 {
         %0 = %1 + %2;
         GOTO bb1;
    }
    bb1 {
            return;
    }
}
```

- %0 is the return value
- Variables are mutable in general
- Basic blocks are identified. We use this to create safe points for yield.
 - We use 'Duff's device' like approach to compile Ballerina to support user space threading
- IR is typed



Ballerina IR Generation

- Simple recursive depth first walk over the graph
- Users visitor pattern
- Pros : simple, local decision making
- Cons: repeated code
 - Fix: post-code gen cleanup.

https://github.com/ballerina-platform/ballerina-lang/blob/master/compiler/ballerina-lang/src/main/jav a/org/wso2/ballerinalang/compiler/bir/BIRGen.java

JVM bytecode Generation

- Iterate BIR and generate JVM code using ASM library
- Ballerina supports user space scheduling. We generate a switch that can jump to any BasicBlock from the top to the function. This helps us resume a function (AKA: Duff's device).
- To support Ballerina's structural typing we lower property access to Map.get() but actual implementation of the "Map" can be a specifically generated class

https://github.com/ballerina-platform/ballerina-lang/blob/master/compiler/ballerina-lang/src/main/jav a/org/wso2/ballerinalang/compiler/bir/codegen/methodgen/MethodGen.java



https://github.com/manuranga/ir-gen-uom

Small task to familiarize yourself with ir generation and Ballerina language. Send us over discord <u>https://discord.com/invite/ballerinalang</u> for a small gift.

